



Dark-Pool Smart Order Routing: a Combinatorial Multi-armed Bandit Approach

Martino Bernasconi, Stefano Martino, **Edoardo Vittori**, Francesco Trovò, Marcello Restelli Smart Order Routing (SOR): optimally splitting an order over multiple venues







Regulated Exchanges - Limit Order Book







Dark Pools







Dark Pool Smart Order Routing - DPSOR

Task	Assumptions	Formulation
 Create and maintain a estimate of hidden liquidity of multiple dark pools Make optimal joint routing and pricing decisions Optimize the dollar volume 	<list-item></list-item>	 Sequential decision problem where at each round t, an agent, given a volume V of shares to execute, must maximize the dollar volume by allocating the shares across K dark pools, specifying the price





Joint routing and pricing allocation







Problem formalization and notation



- A_{kn}^t : amount allocated at round t to dark pool k at price p_n
 - We have the constraint that

$$\sum_{k=1}^{K} \sum_{n=1}^{N} A_{kn}^{t} = V_{kn}$$

• Our objective is the allocation that maximizes dollar volume

$$R_{t}(\mathfrak{U}) = \sum_{k=1}^{K} \sum_{n=1}^{N} r_{kn}^{t} p_{n}$$
Censored feedback
$$r_{kn}^{t} = \min\{A_{kn}^{t}, s_{kn}^{t}\}$$

 s_{kn}^t is the actual liquidity present at time t in dark pool k at price p_n





Censored feedback



Volume



POLITECNICO MILANO 1863



Combinatorial MAB [Chen et al., 2013]



• We are in a CMAB setting, where the superarms are all the combinations of A_{kn}^t which satisfy the following constraint:

$$\sum_{k=1}^{K} \sum_{n=1}^{N} A_{kn}^{t} = V$$

• We want to minimize pseudo-regret w.r.t. the expected dollar value of the optimal superarm r^*

$$Reg_{t}(\mathfrak{U}) := t r^{*} - \sum_{h=1}^{t} \sum_{k=1}^{K} \sum_{n=1}^{N} \mathbb{E}[r_{kn}^{h}] \mathbb{1}\{A_{nk}^{h} > 0\} p_{n}$$
$$r_{kn}^{t} = \min\{A_{kn}^{t}, s_{kn}^{t}\}$$





Estimating liquidity



Dark Pool k





Counting successes $oldsymbol{lpha}$ and failures $oldsymbol{eta}$







DP-CMAB Algorithm – θ Selection



Translating liquidity to allocation







At each round t:

- Calculate the liquidity estimate θ_t using α_t , β_t and the appropriate update Bayes, CUCB or TS
- Calculate the action matrix $A_t \leftarrow \text{Opt}(\theta_t)$
- Play allocation A_t
- Receive feedbacks r_t from played arms
- Calculate the parameters $lpha_{t+1}$ and eta_{t+1}











Experimental setup

- Selling V = 10 units
- K = 10 dark pools
- Prices in {90, 91, ..., 100}
- Rounds T = 1000
- Results averaged over 20 runs







Experimental results – Different strategies and propagations







Related works

Censored Exploration and the Dark Pool Problem, Ganchev et al. [2009]

• Finds an optimal allocation strategy based on the Kaplan-Meier estimator

Optimal Allocation Strategies for the Dark Pool Problem, Agarwal [2010]

• Extends Ganchev et al. [2009] to an adversarial scenario

Limitations

They do not allow the agent to specify the transaction price and do not take advantage of domain knowledge



Price selection

- Random: selects a random price at the beginning of each round
- Oracle: selects best single price which maximizes the expected cumulated dollar volume across the run.





Experimental results – Comparison with baselines







Experimental results – Dollar volume







Conclusions

We developed the DP-CMAB algorithm that:

- extends the limitations of state-of-the-art algorithms
- exploits the knowledge about the financial setting to improve its allocation policy
- empirically outperforms state-of-the-art algorithms in a realistic scenario

Further extensions:



- regret proof of propagation updates
- nonstationary (ADVERSARIAL) dark pool liquidity









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